

THE NUMBER OF PARKING FUNCTIONS WITH CENTER OF A GIVEN LENGTH

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ABSTRACT. Let $1 \leq r \leq n$ and suppose that, when the *Depth-first Search Algorithm* is applied to a given rooted labeled tree on $n + 1$ vertices, exactly r vertices are visited before backtracking. Let R be the set of trees with this property. We count the number of elements of R .

For this purpose, we first consider a bijection, due to Perkinson, Yang and Yu, that maps R onto the set of parking function with *center* (defined by the authors in a previous article) of size r . A second bijection maps this set onto the set of parking functions with *run* r , a property that we introduce here. We then prove that the number of length n parking functions with a given run is the number of length n *rook words* (defined by Leven, Rhoades and Wilson) with the same run. This is done by counting related lattice paths in a ladder-shaped region. We finally count the number of length n *rook words* with run r , which is the answer to our initial question.

1. INTRODUCTION

Let \mathbb{T}_n be the set of rooted labeled trees on the set of vertices $\{0, 1, \dots, n\}$ with root $r = 0$, and let $T \in \mathbb{T}_n$. Suppose that the *Depth-first Search Algorithm* (DFS) is applied to T by starting at r and by visiting at each vertex the (unvisited) neighbor of highest label. If T is not a path with endpoint r , at a certain moment the algorithm will backtrack. In this paper we are concerned with the number of vertices that are visited before this happens.

More precisely, let $\mathbf{v} = \mathbf{v}(T) = (v_1, \dots, v_k)$ be the ordered set of vertices different from the root that are visited *before backtracking*, and let $\text{leg}(T) = k$ be the length of \mathbf{v} . We evaluate here explicitly the enumerator

$$\mathcal{L}\mathcal{L}\mathcal{T}_n(t) = \sum_{T \in \mathbb{T}_n} t^{\text{leg}(T)}$$

It is well-known that $|\mathbb{T}_n| = (n + 1)^{n-1} = |\text{PF}_n|$, where PF_n is the set of *parking functions of length* n , consisting of the n -tuples $\mathbf{a} = (a_1, \dots, a_n)$ of positive integers such that the i th entry *in ascending order* is always at most $i \in [n] := \{1, \dots, n\}$. As usual, we may denote \mathbf{a} either as a word, $\mathbf{a} = a_1 \cdots a_n$, or as a function, $\mathbf{a}: [n] \rightarrow [n]$ such that $\mathbf{a}(i) = a_i$.

Parking functions are an important combinatorial structure with several connections to other areas of mathematics (see e.g. Haglund [6] and the recent survey by Yan [15]). In particular, following Kreweras [9], various bijections between trees on $n + 1$ vertices and parking functions of length n were defined where the reversed sum enumerator for parking functions is the counterpart to the inversion enumerator for trees [5, 7, 12, 15].

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In a recent paper, Perkinson, Yang and Yu [11] constructed a very general algorithm that gives us as a particular case a new bijection with this property.

We show in Section 2 (cf. [3]) that, for this bijection, the counterpart of the statistics $\text{leg}(T)$ is $z(\mathbf{a}) = |Z(\mathbf{a})|$, where $Z(\mathbf{a})$ is *the center of \mathbf{a}* defined in [2]. We recall that, for any $\mathbf{a} \in [n]^n$, $X = \{x_1, \dots, x_\ell\}$ with $1 \leq x_1 \leq \dots \leq x_\ell \leq n$ is *\mathbf{a} -central* if

$$a_{x_i} \leq i \quad (i = 1, \dots, \ell)$$

Clearly, if X and Y are \mathbf{a} -central so it is $X \cup Y$. The center of \mathbf{a} is the (unique) *maximal \mathbf{a} -central subset $Z(\mathbf{a})$ of $[n]$* . We prove namely that if $\mathbf{v}(T) = (v_1, \dots, v_k)$ then exactly $Z(\mathbf{a}) = \{v_1, \dots, v_k\}$. Hence, we obtain, if we consider the enumerator

$$\mathcal{ZPF}_n(t) = \sum_{\mathbf{a} \in \text{PF}_n} t^{z(\mathbf{a})},$$

Theorem 2.1. *For every $n \in \mathbb{N}$,*

$$\mathcal{LLT}_n(t) = \mathcal{ZPF}_n(t).$$

The evaluation of this new enumerator was indeed part of our initial twofold purpose for its role in the theory of parking functions, described as follows. Consider the Shi arrangement, formed by all the hyperplanes defined in \mathbb{R}^n by equations of form $x_i - x_j = 0$ and of form $x_i - x_j = 1$, where $1 \leq i < j \leq n$. Let R_0 be the chamber of the arrangement consisting of the intersection of all the open slabs defined by the condition $0 < x_i - x_j < 1$. Pak and Stanley [13] defined a bijective labeling of the chambers of this arrangement by parking functions, in which R_0 is labeled with the parking function $(1, \dots, 1)$ and, along a shortest path from R_0 to any other chamber, for any crossed hyperplane of form $x_i - x_j = 0$ the j th coordinate of the label is increased by one, and for any crossed plane of form $x_i - x_j = 1$ it is the i th coordinate that is increased by one. The bijection is defined from chambers (represented by permutations of $[n]$ decorated with arcs following certain rules) to parking functions. See [4] for a very general perspective of this bijection. *Centers* occur in the opposite direction, i.e., it is from the *center* of any parking function that we may recover the chamber that it labels in the Pak-Stanley labeling [2].

For example, consider the region of the Shi arrangement in \mathbb{R}^9 defined by

$$\begin{aligned} x_8 < x_4 < x_3 < x_9 < x_6 < x_7 < x_1 < x_2 < x_5, \\ x_8 + 1 > x_7, x_3 + 1 > x_2, x_7 + 1 > x_5, \\ x_4 + 1 < x_1, x_6 + 1 < x_5. \end{aligned}$$

This region may be represented by

$$\overbrace{843967125}.$$

and is associated by the Pak-Stanley bijection to the parking function of Example 2.2,

$$\mathbf{a} = 341183414.$$

The center of the parking function determines the first elements of the permutation, which include the elements under the first arc, in the right order, 843967 (cf. Example 2.2). By knowing this, we may replace the parking function by another one of shorter length, and proceed recursively [2].

We now consider a third statistic. Let, for $\mathbf{a} = (a_1, \dots, a_n) \in [n]^n$ such that $1 \in \{a_1, \dots, a_n\}$,

$$\text{run}(\mathbf{a}) = \max \{i \in [n] \mid [i] \subseteq \{a_1, \dots, a_n\}\},$$

and let $\text{run}(\mathbf{a}) = 0$ if $1 \notin \{a_1, \dots, a_n\}$. We prove in Section 3 the following result. Let

$$\mathcal{RPF}_n(t) = \sum_{\mathbf{a} \in \text{PF}_n} t^{\text{run}(\mathbf{a})}.$$

Theorem 3.5. *For every $n \in \mathbb{N}$,*

$$\mathcal{ZPF}_n(t) = \mathcal{RPF}_n(t).$$

Now, consider the set RW_n of *rook words of length n* defined by Leven, Rhoades and Wilson [10], that is, the ordered sets $\mathbf{a} = (a_1, \dots, a_n) \in [n]^n$ such that $a_1 \leq \text{run}(\mathbf{a})$. Let

$$\mathcal{RRW}_n(t) = \sum_{\mathbf{a} \in \text{RW}_n} t^{\text{run}(\mathbf{a})}$$

The key to our enumeration is developed in Section 4, where we prove the following result.

Theorem 4.11. *For every $n \in \mathbb{N}$,*

$$\mathcal{RPF}_n(t) = \mathcal{RRW}_n(t).$$

In this case we do not consider all parking functions at once. Instead, we only consider those for which the sets of elements with the same image are the same. We count parking functions by counting nonnegative sequences that are componentwise bounded above by a given positive sequence. Based on results of independent interest we prove that their number is the number of rook words defined in the same way.

Finally, by directly counting rook words with a given run, we are able to present in Section 5 an expression for all the (equal) previous enumerators.

Theorem 5.1. *For every integers $1 \leq r \leq n$,*

$$\begin{aligned} [t^r](\mathcal{L}\mathcal{L}\mathcal{T}_n(t)) &= r! \sum_{i_1 + \dots + i_r = n-r} (n-1)^{i_1} (n-2)^{i_2} \dots (n-r)^{i_r} \\ &= r \sum_{j=0}^{r-1} (-1)^j \binom{r-1}{j} (n-1-j)^{n-1}. \end{aligned}$$

It is perhaps worth noting that rook words were introduced in order to label the chambers of the *Ish arrangement*, defined in \mathbb{R}^n by all the hyperplanes with equations of form $x_i - x_j = 0$, as before, and of form $x_1 - x_j = i$, where again $1 \leq i < j \leq n$. Several bijections, which preserve different properties, have been defined between the chambers of the Shi arrangement and the chambers of the Ish arrangement, particularly by Leven, Rhoades and Wilson using rook words [10]. In fact, our work here may be presented as another example of a general statement by Armstrong and Rhoades [1], saying that “The Ish arrangement is something of a ‘toy model’ for the Shi arrangement”, in the sense that several properties are shared by both arrangements, but are easier to prove in case of the Ish arrangement than in the case of the Shi arrangement.

Example 1.1. We consider in Table 1 the case where $n = 3$ and hence

$$\mathcal{L}\mathcal{L}\mathcal{T}_3(t) = \mathcal{ZPF}_3(t) = \mathcal{RPF}_3(t) = \mathcal{RRW}_3(t) = 4t + 6t^2 + 6t^3$$

by classifying the corresponding trees and parking functions according to the various statistics and the corresponding bijections. Note that for $n = 3$ rook words and parking functions coincide, except that $311 \in \text{PF}_3 \setminus \text{RW}_3$ whereas $133 \in \text{RW}_3 \setminus \text{PF}_3$. But $\text{run}(311) = \text{run}(133) = 1$.

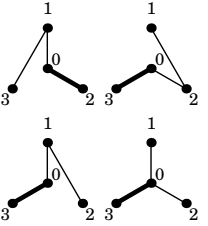
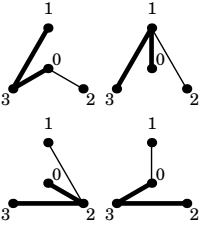
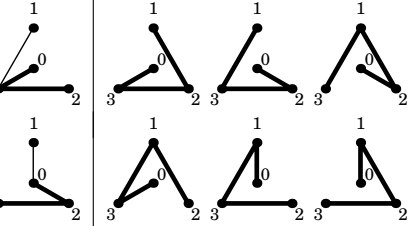
k	1	2	3
Trees T with $\text{leg}(T) = k$			
Parking functions \mathbf{a} with $z(\mathbf{a}) = k$	$\underline{213}$ $\underline{221}$ $\underline{231}$ $\underline{321}$	$\underline{131}$ $\underline{132}$ $\underline{211}$ $\underline{212}$ $\underline{311}$ $\underline{312}$	$\underline{111}$ $\underline{112}$ $\underline{113}$ $\underline{121}$ $\underline{122}$ $\underline{123}$
Parking functions \mathbf{a} with $\text{run}(\mathbf{a}) = k$	$\underline{113}$ $\underline{111}$ $\underline{131}$ $\underline{311}$	$\underline{211}$ $\underline{122}$ $\underline{221}$ $\underline{112}$ $\underline{121}$ $\underline{212}$	$\underline{321}$ $\underline{231}$ $\underline{213}$ $\underline{312}$ $\underline{132}$ $\underline{123}$

TABLE 1. The case where $n = 3$ 2. FROM LABELED TREES TO PARKING FUNCTIONS: LEGS *vs.* CENTERS

We reproduce here the algorithm (Algorithm 1, below) of Perkinson, Yang and Yu [11], that takes as input the parking function \mathbf{a} (or, more precisely, takes as input $\mathcal{P} = \mathbf{a} - 1: [n] \rightarrow \mathbb{Z}$ such that $\mathcal{P}(i) = a_i - 1$) and returns the list `tree_edges` of edges of a tree. Note that a parking function is a G -parking function where G is a (rooted) complete graph on V . For this reason, the sentence “for each j adjacent to i in G ” reads as “for each $j \neq i$ ”.

Note also that, in general, a spanning tree T of G is seen as a directed graph in which all paths lead away from the root. So, edge ij is written (i, j) if i is in the (unique) path from i to j (with no vertex between them). Note also that, by definition (cf. Line 7), if after running the algorithm both edges (i, j) and (i, k) belong to T and $j > k$ then `DFS_BURN(j)` has been called before `DFS_BURN(k)`.

Theorem 2.1. *For every $n \in \mathbb{N}$,*

$$\mathcal{L}\mathcal{L}\mathcal{T}_n(t) = \mathcal{Z}\mathcal{P}\mathcal{F}_n(t)$$

Proof. We show that there exist a bijection $\varphi: \text{PF}_n \rightarrow \mathbb{T}_n$ such that, for every $\mathbf{a} \in \text{PF}_n$, if $T = \varphi(\mathbf{a}) \in \mathbb{T}_n$, then $\text{leg}(T) = z(\mathbf{a})$.

Let T be the tree given by Algorithm 1 with input $\mathcal{P} = \mathbf{a} - 1$ (we know this defines a bijection from PF_n to \mathbb{T}_n by [11, Theorem 3]). Now, let ℓ be the first value of i where, when `DFS_FROM(i)` is called, $\mathcal{P}(j) > 0$ whenever $j \notin \text{burnt_vertices}$. If this never occurs, let ℓ be the last vertex joined to `burnt_vertices`.

Let $B = (0 = v_0, \dots, v_k = i) = \text{burnt_vertices}$ and $E = \text{tree_edges}$ at the end of the loop of `DFS_FROM(i)` (the end of Line 14) for $i = \ell$, and note that, by definition, $E = ((v_0, v_1), \dots, (v_{k-1}, v_k))$. Hence, $\mathbf{v}(T) = (v_1, \dots, v_k)$.

Now, let $X = \{x_1, \dots, x_k\} = \{v_1, \dots, v_k\}$ with $x_1 < \dots < x_k$.

We must prove that:

- (1) for every $m \in [k]$, $\mathbf{a}(x_m) \leq m$;
- (2) X is maximal for this property;

Clearly, if $x_{i_1} = v_1$, then $\mathbf{a}(x_{i_1}) \leq i_1$ since, by definition, $v_1 = \max(\mathcal{P}^{-1}(\{0\}))$ and so $\mathbf{a}(x_{i_1}) = 1$. Now, suppose that the same holds true for $x_{i_2} = v_2, \dots, x_{i_{\ell-1}} = v_{\ell-1}$, consider $x_{i_\ell} = v_\ell$ and note that, when `DFS_FROM($v_{\ell-1}$)` is called, v_ℓ is the largest value

Algorithm 1 DFS-burning algorithm.

ALGORITHM

Input: $\mathcal{P}: V \setminus \{r\} \rightarrow \mathbb{N} \cup \{0\}$ 1: `burnt_vertices` = $\{r\}$ 2: `dampened_edges` = $\{\}$ 3: `tree_edges` = $\{\}$ 4: execute `DFS_FROM`(r)**Output:** `burnt_vertices` and `tree_edges`

AUXILLARY FUNCTION

5: **function** `DFS_FROM`(i)6: **foreach** j adjacent to i in G , from largest numerical value to smallest **do**7: **if** $j \notin$ `burnt_vertices` **then**8: **if** $\mathcal{P}(j) = 0$ **then**9: append j to `burnt_vertices`10: append (i, j) to `tree_edges`11: `DFS_FROM`(j)12: **else**13: $\mathcal{P}(j) = \mathcal{P}(j) - 1$ 14: append (i, j) to `dampened_edges`

of $j \notin \{v_0, v_1, \dots, v_{\ell-1}\}$ with $\mathcal{P}(j) = 0$. Since $\mathcal{P}(v_\ell)$ has been reduced in earlier calls to `DFS_FROM`(v_m) (in Line 15) exactly when $v_m < v_\ell$, since it is now zero, and since new additions to `burnt_vertices` will not decrease the order of v_ℓ in the corresponding set, $\mathbf{a}(x_{i_\ell}) \leq i_\ell$.

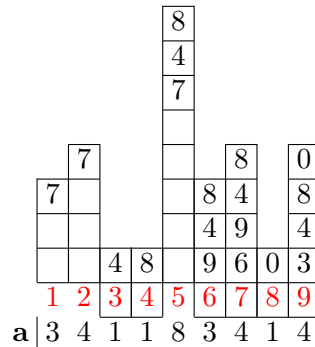
When finally `DFS_FROM`(v_k) is called, $\mathcal{P}(j) > 0$ for all $j \notin X$. In particular, for each such j , the number of elements in X less than or equal to j is necessarily less than a_j . \square

More precisely, if $\mathbf{v}(T) = (v_1, \dots, v_k)$, then clearly

$$\mathbf{a}(v_i) = |\{k \in [i] \mid v_k \leq v_i\}|.$$

Compare with Definition 3.1 below.

Example 2.2. Let $\mathbf{a} = 341183414 \in \text{PF}_9$. We apply Algorithm 1 to \mathbf{a} by drawing a_j empty boxes for each $j \in [9]$ that are filled with i during the execution of `DFS_FROM`(i), at Line 14. Below, `DFS_FROM`(i) has been called for, in this order, $i = 0, 8, 4, 3, 9, 6, 7$. Hence, at the moment, $i = 7$ and `burnt_vertices` = $(0, 8, 4, 3, 9, 6, 7)$. Since $\mathcal{P}(j) > 0$ for $j \notin$ `burnt_vertices` (i.e., for $j = 1, 2, 5$), $\ell = i = 7$, and so $\mathbf{v} = (8, 4, 3, 9, 6, 7)$.



3. WITHIN PARKING FUNCTIONS: CENTERS *vs.* RUNS

Definition 3.1. Consider, for a positive integer n and for a permutation $\mathbf{w} = (w_1, \dots, w_n) \in \mathfrak{S}_n$,

$$f_{w_i} = |\{k \in [i] \mid w_k \leq w_i\}|, \quad i = 1, \dots, n,$$

and

$$t_n(\mathbf{w}) = (f_1, \dots, f_n) \in [1] \times [2] \times \dots \times [n].$$

According to [2], t_n is a bijection between \mathfrak{S}_n and $[1] \times [2] \times \dots \times [n]$.

Example 3.2. If $\mathbf{w} = 521634$, then $f_1 = f_{w_3} = 1$, $f_2 = f_{w_2} = 1$, $f_3 = f_{w_5} = 3$, $f_4 = f_{w_6} = 4$, $f_5 = f_{w_1} = 1$ and $f_6 = f_{w_4} = 4$. Hence $t(\mathbf{w}) = 113414 \in [1] \times \dots \times [6]$.

Given $\mathbf{a} \in [n]^n$, let

$$\text{Run}(\mathbf{a}) = \{\max \mathbf{a}^{-1}(\{j\}) \mid 1 \leq j \leq \text{run}(\mathbf{a})\}$$

if $\text{run}(\mathbf{a}) > 0$, and let $\text{Run}(\mathbf{a}) = \emptyset$ if $\text{run}(\mathbf{a}) = 0$. Then $|\text{Run}(\mathbf{a})| = \text{run}(\mathbf{a})$.

For $A \subseteq [n]$, let

$$Z_n^{-1}(A) = \{\mathbf{a} \in [n]^n \mid Z(\mathbf{a}) = A\}$$

and

$$\text{Run}_n^{-1}(A) = \{\mathbf{a} \in [n]^n \mid \text{Run}(\mathbf{a}) = A\}.$$

Now let $A = \{i_1, i_2, \dots, i_k\} \neq \emptyset$ with $i_1 < i_2 < \dots < i_k$ and take $i_0 = 0$ and $i_{k+1} = n + 1$. If $\mathbf{a} = (a_1, a_2, \dots, a_n) \in Z_n^{-1}(A)$, then

- $(a_{i_1}, a_{i_2}, \dots, a_{i_k}) \in [1] \times [2] \times \dots \times [k]$,
- $a_j \in \{\ell + 1, \dots, n\} = [n] \setminus [\ell]$, if $i_{\ell-1} < j < i_\ell$, with $\ell \in [k + 1]$.

$$\begin{array}{ccccccc} \neq 1 & \boxed{a_{i_1} \leq 1} & \neq \begin{Bmatrix} 1 \\ 2 \end{Bmatrix} & \boxed{a_{i_2} \leq 2} & \neq \begin{Bmatrix} 1 \\ 2 \\ 3 \end{Bmatrix} & \dots & \neq \begin{Bmatrix} 1 \\ \vdots \\ k \end{Bmatrix} & \boxed{a_{i_k} \leq k} & \neq \begin{Bmatrix} 1 \\ \vdots \\ k+1 \end{Bmatrix} \\ \hline & \underbrace{\hspace{1.5cm}}_{i_1} & & \underbrace{\hspace{1.5cm}}_{i_2} & & \dots & & \underbrace{\hspace{1.5cm}}_{i_k} & \end{array}$$

If $\mathbf{a} = (a_1, a_2, \dots, a_n) \in \text{Run}_n^{-1}(A)$, then

- $(a_{i_1}, a_{i_2}, \dots, a_{i_k}) \in \mathfrak{S}_k$,
- $a_j \in [n] \setminus \{k + 1, a_{i_1}, \dots, a_{i_{\ell-1}}\}$, if $i_{\ell-1} < j < i_\ell$, with $\ell \in [k + 1]$.

$$\begin{array}{ccccccc} \neq k+1 & \boxed{a_{i_1}} & \neq \begin{Bmatrix} k+1 \\ a_{i_1} \end{Bmatrix} & \boxed{a_{i_2}} & \neq \begin{Bmatrix} k+1 \\ a_{i_1} \\ a_{i_2} \end{Bmatrix} & \dots & \neq \begin{Bmatrix} k+1 \\ \vdots \\ a_{i_{k-1}} \end{Bmatrix} & \boxed{a_{i_k}} & \neq \begin{Bmatrix} k+1 \\ \vdots \\ a_{i_k} \end{Bmatrix} \\ \hline & \underbrace{\hspace{1.5cm}}_{i_1} & & \underbrace{\hspace{1.5cm}}_{i_2} & & \dots & & \underbrace{\hspace{1.5cm}}_{i_k} & \end{array}$$

Clearly, the size of both $Z_n^{-1}(A)$ and $\text{Run}_n^{-1}(A)$ is

$$k!(n-1)^{i_1-1}(n-2)^{i_2-1} \dots (n-k)^{i_k-1}$$

if $|A| = k > 0$, and $(n-1)^n$ if $A = \emptyset$.

We now define two mappings $\Phi, \Psi : [n]^n \rightarrow [n]^n$.

If $Z(\mathbf{a}) = \emptyset$, we define $\Phi(\mathbf{a}) := \mathbf{a}$. Else, if $Z(\mathbf{a}) = \{i_1, i_2, \dots, i_k\} \neq \emptyset$ with $i_1 < i_2 < \dots < i_k$, we define $\Phi(\mathbf{a})$ as follows. Let $\mathbf{b} := b_1 b_2 \dots b_k = t_k^{-1}(a_{i_1} a_{i_2} \dots a_{i_k}) \in \mathfrak{S}_k$ and

$\sigma_{\mathbf{a}} \in \mathfrak{S}_n$ be the permutation of length n defined by

$$\sigma_{\mathbf{a}}(j) = \begin{cases} k+1 & \text{if } j = 1; \\ b_{j-1} & \text{if } 2 \leq j \leq k+1; \\ j & \text{if } k+2 \leq j \leq n. \end{cases}$$

Finally, let

$$(\Phi(\mathbf{a}))(j) := \begin{cases} b_\ell & \text{if } j = i_\ell \in Z(\mathbf{a}); \\ \sigma_{\mathbf{a}}(a_j) & \text{if } j \notin Z(\mathbf{a}). \end{cases}$$

If $\text{Run}(\mathbf{a}) = \emptyset$, we define $\Psi(\mathbf{a}) := \mathbf{a}$. Else, if $\text{Run}(\mathbf{a}) = \{i_1, i_2, \dots, i_k\} \neq \emptyset$ with $i_1 < i_2 < \dots < i_k$, we define $\Psi(\mathbf{a})$ as follows. Let $\mathbf{c} := c_1 c_2 \dots c_k = t_k(a_{i_1} a_{i_2} \dots a_{i_k}) \in [1] \times [2] \times \dots \times [k]$ and $\tau_{\mathbf{a}} \in \mathfrak{S}_n$ be the permutation of length n defined by

$$\tau_{\mathbf{a}}(j) = \begin{cases} c_j + 1 & \text{if } 1 \leq j \leq k; \\ 1 & \text{if } j = k+1; \\ j & \text{if } k+2 \leq j \leq n. \end{cases}$$

Finally, let

$$(\Psi(\mathbf{a}))(j) := \begin{cases} c_\ell & \text{if } j = i_\ell \in \text{Run}(\mathbf{a}); \\ \tau_{\mathbf{a}}(a_j) & \text{if } j \notin \text{Run}(\mathbf{a}). \end{cases}$$

Example 3.3. Let $\mathbf{a} = 341183414 \in [9]^9$. On the one hand, $Z(\mathbf{a}) = \{3, 4, 6, 7, 8, 9\}$, $t_6^{-1}(a_3 a_4 a_6 a_7 a_8 a_9) = t_6^{-1}(113414) = 521634 \in \mathfrak{S}_6$, and $\sigma_{\mathbf{a}} = \underline{752163489}$, so $\Phi(\mathbf{a}) = \underline{215281634}$.

$$\mathbf{a} = \begin{array}{cccccccc} 3 & 4 & \boxed{1} & \boxed{1} & 8 & \boxed{3} & \boxed{4} & \boxed{1} & \boxed{4} \\ & & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} \\ & & 3 & 4 & & 6 & 7 & 8 & 9 \end{array}$$

$$\Phi(\mathbf{a}) = \begin{array}{cccccccc} 2 & 1 & \boxed{5} & \boxed{2} & 8 & \boxed{1} & \boxed{6} & \boxed{3} & \boxed{4} \\ & & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} \\ & & 3 & 4 & & 6 & 7 & 8 & 9 \end{array}$$

On the other hand, $\text{Run}(\mathbf{a}) = \{8\}$, $\mathbf{c} = t_1(a_8) = t_1(1) = 1$ and $\tau_{\mathbf{a}} = \underline{213456789}$, so $\Psi(\mathbf{a}) = \underline{342283414}$. Note that \mathbf{a} belongs to PF_9 , as well as $\Phi(\mathbf{a})$ and $\Psi(\mathbf{a})$.

The following theorem summarizes the main properties of Φ and Ψ .

Theorem 3.4.

- (1) For all $\mathbf{a} \in [n]^n$, $z(\mathbf{a}) = \text{run}(\Phi(\mathbf{a}))$, $Z(\mathbf{a}) = \text{Run}(\Phi(\mathbf{a}))$, $\text{run}(\mathbf{a}) = z(\Psi(\mathbf{a}))$, and $\text{Run}(\mathbf{a}) = Z(\Psi(\mathbf{a}))$,
- (2) Φ and Ψ are bijections and $\Psi = \Phi^{-1}$,
- (3) $\Phi(\text{PF}_n) = \text{PF}_n$.

Proof.

(1) Let $Z(\mathbf{a}) = \{i_1, \dots, i_k\}$ with $i_1 < \dots < i_k$ and $\Phi(\mathbf{a}) =: \mathbf{d} = d_1 \dots d_n$. On the one hand, $k \leq \text{run}(\mathbf{d})$ since $\{d_{z_1}, \dots, d_{z_k}\} = [k]$. On the other hand, $k+1 \notin \{d_1, \dots, d_n\}$ because $k+1 \notin \{d_{z_1}, \dots, d_{z_k}\}$ and there is no $j \in [n] \setminus Z(\mathbf{a})$ such that $\mathbf{a}(j) = 1$. Hence $z(\mathbf{a}) = k = \text{run}(\mathbf{d})$.

Let $i_j \in Z(\mathbf{a})$. First, note that $d_{i_j} \leq j \leq k$. Second, note that if $\ell > i_j$, then $d_\ell \neq d_{i_j}$; if $\ell = i_p \in Z(\mathbf{a}) \setminus \{i_j\}$, then $d_\ell \neq d_{i_j}$ because $d_{z_1} \dots d_{z_k} \in \mathfrak{S}_k$; if $\ell \in [n] \setminus (\{i_j\} \cup Z(\mathbf{a}))$, then $\mathbf{a}(\ell) > j+2$ and $d_\ell \neq d_{i_j}$, since $\sigma_{\mathbf{a}} \in \mathfrak{S}_n$.

Similarly, one can show that $\text{run}(\mathbf{a}) = z(\Psi(\mathbf{a}))$ and $\text{Run}(\mathbf{a}) = Z(\Psi(\mathbf{a}))$.

(2) Given $\mathbf{a} \in [n]^n$, we have $Z(\mathbf{a}) = \text{Run}(\Phi(\mathbf{a}))$, $\text{Run}(\mathbf{a}) = Z(\Psi(\mathbf{a}))$, $\tau_{\Phi(\mathbf{a})} = \sigma_{\mathbf{a}}^{-1}$ and $\sigma_{\Psi(\mathbf{a})} = \tau_{\mathbf{a}}^{-1}$. Hence $(\Psi \circ \Phi)(\mathbf{a}) = \mathbf{a} = (\Phi \circ \Psi)(\mathbf{a})$.

(3) Let $\mathbf{a} \in \text{PF}_n$ and $k = z(\mathbf{a})$. If $j \leq k$, $|\Phi(\mathbf{a})^{-1}([j])| \geq j$, since $[j] \subseteq [k] \subseteq \Phi(\mathbf{a})([n])$. If $j > k$, then $\Phi(\mathbf{a})^{-1}([j]) = \mathbf{a}^{-1}([j])$ and so $|\Phi(\mathbf{a})^{-1}([j])| = |\mathbf{a}^{-1}([j])| \geq j$ because $\mathbf{a} \in \text{PF}_n$. Since $\Phi(\text{PF}_n) \subseteq \text{PF}_n$ and Φ is a bijection, $\Phi(\text{PF}_n) = \text{PF}_n$. \square

Theorem 3.5. *For every $n \in \mathbb{N}$,*

$$\mathcal{ZPF}_n(t) = \mathcal{RPF}_n(t).$$

4. FROM PARKING FUNCTIONS TO ROOK WORDS

4.1. Restricted integer sequences. We start this section by considering a general situation of independent interest.

Definition 4.1. Let, for a positive integer k and for $\boldsymbol{\ell} = (\ell_1, \dots, \ell_k) \in \mathbb{N}^k$, $\mathbf{L} = (L_1, \dots, L_k) \in \mathbb{N}^k$ be the cumulative sum of $\boldsymbol{\ell}$, i.e.,

$$L_i = \ell_1 + \ell_2 + \dots + \ell_i, \quad i = 1, \dots, k,$$

and consider the set

$$\langle \ell_1, \dots, \ell_k \rangle = \{(x_0, x_1, \dots, x_k) \in \mathbb{Z}^{k+1} \mid x_0 = 0; \forall 1 \leq i \leq k, x_{i-1} < x_i \leq L_i\}$$

Lemma 4.2. *For every positive integers k, ℓ_1, \dots, ℓ_k , if $i < k$ and $\ell_{i+1} > 1$, then*

$$\begin{aligned} |\langle \ell_1, \dots, \ell_{i-1}, \ell_i + 1, \ell_{i+1} - 1, \ell_{i+2}, \dots, \ell_k \rangle| \\ = |\langle \ell_1, \dots, \ell_k \rangle| + |\langle \ell_1, \dots, \ell_{i-1} \rangle| |\langle \ell_{i+1} - 1, \ell_{i+2}, \dots, \ell_k \rangle| \end{aligned}$$

whereas

$$|\langle \ell_1, \dots, \ell_{k-1}, \ell_k + 1 \rangle| = |\langle \ell_1, \dots, \ell_k \rangle| + |\langle \ell_1, \dots, \ell_{k-1} \rangle|$$

and, if $\ell_1 > 1$,

$$|\langle \ell_1 - 1, \ell_2, \dots, \ell_k \rangle| = |\langle \ell_1, \dots, \ell_k \rangle| - |\langle \ell_1 + \ell_2 - 1, \ell_3, \dots, \ell_k \rangle|.$$

Proof. We present here a bijective proof. Note that, for every $i < k$,

$$\langle \ell_1, \dots, \ell_k \rangle \subseteq \langle \ell_1, \dots, \ell_{i-1}, \ell_i + 1, \ell_{i+1} - 1, \ell_{i+2}, \dots, \ell_k \rangle.$$

But, by definition,

$$\begin{aligned} (x_0, \dots, x_k) \in \langle \ell_1, \dots, \ell_{i-1}, \ell_i + 1, \ell_{i+1} - 1, \ell_{i+2}, \dots, \ell_k \rangle \setminus \langle \ell_1, \dots, \ell_k \rangle \\ \iff \begin{cases} x_0 = 0 \\ x_i = L_i + 1 \\ x_{j-1} < x_j \leq L_j \text{ for every } j \neq i \text{ with } 1 \leq j \leq k, \end{cases} \\ \iff \begin{cases} (x_0, \dots, x_{i-1}) \in \langle \ell_1, \dots, \ell_{i-1} \rangle \\ (x_i - L_i - 1, \dots, x_k - L_i - 1) \in \langle \ell_{i+1}, \dots, \ell_k \rangle \end{cases} \end{aligned}$$

For the second statement, note that also $\langle \ell_1, \dots, \ell_k \rangle \subseteq \langle \ell_1, \dots, \ell_{k-1}, \ell_k + 1 \rangle$ and that $(x_0, \dots, x_k) \in \langle \ell_1, \dots, \ell_{k-1}, \ell_k + 1 \rangle \setminus \langle \ell_1, \dots, \ell_k \rangle$ if and only if $x_k = \ell_k + 1$ and $(x_0, \dots, x_{k-1}) \in \langle \ell_1, \dots, \ell_{k-1} \rangle$.

Finally, for the third statement, note that, by definition, if

$$(0, x_1, \dots, x_k) \in \langle \ell_1 - 1, \ell_2, \dots, \ell_k \rangle,$$

then

$$\begin{cases} x_1 + 1 > 1 \\ (0, x_1 + 1, \dots, x_k + 1) \in \langle \ell_1, \ell_2, \dots, \ell_k \rangle. \end{cases}$$

In fact, given a k -tuple $(y_1, \dots, y_k) \in \mathbb{N}^k$,

$$\begin{cases} (0, y_1, \dots, y_k) \in \langle \ell_1, \ell_2, \dots, \ell_k \rangle \\ (0, y_1 - 1, \dots, y_k - 1) \notin \langle \ell_1 - 1, \ell_2, \dots, \ell_k \rangle \end{cases}$$

if and only if

$$\begin{cases} y_1 = 1 \\ (0, y_2 - 1, \dots, y_k - 1) \in \langle \ell_1 + \ell_2 - 1, \ell_3, \dots, \ell_k \rangle. \quad \square \end{cases}$$

Remark 4.3. Let, for any $\mathbf{x} = (0, x_1, \dots, x_k) \in \mathbb{Z}^{k+1}$, $\mathbf{y} = (y_1, \dots, y_k) = (x_1 - 1, x_2 - 2, \dots, x_k - k)$. Then $\mathbf{x} \in \langle \ell_1, \dots, \ell_k \rangle$ if and only if $0 \leq y_1 \leq L_1 - 1$ and $y_i \leq y_{i+1} \leq L_{i+1} - (i + 1)$ for every $i = 1, 2, \dots, k - 1$.

Hence, if (ℓ_1, \dots, ℓ_k) is a *composition of n* (i.e., $n = L_k$) we may represent the elements of $\langle \ell_1, \dots, \ell_k \rangle$ by lattice paths from $(0, 0)$ to $(k, n - k)$ that are contained in the region between the x axis and the path P that has the same ends and the property that the height of the i th horizontal step is $L_i - i$ for every $i = 1, 2, \dots, k$. See Figure 1 for an example. Hence,

$$(4.3.1) \quad |\langle \ell_1, \dots, \ell_k \rangle| = \det_{1 \leq i, j \leq k} \left(\binom{\ell_1 + \dots + \ell_i - i + 1}{j - i + 1} \right).$$

follows (cf. [8, Theorem 10.7.1]). Note that Lemma 4.2 may easily be proved by using the characteristic properties of determinants.

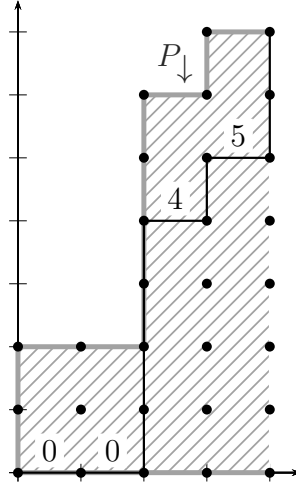


FIGURE 1. Lattice path representation of $(0, 1, 2, 7, 9) \in \langle 3, 1, 5, 2 \rangle$.

Example 4.4. Note that $(0, 1, 2)$ is a subsequence of $\mathbf{x} = (0, 1, 2, 7, 9)$ but $(0, 1, 2, 3)$ is not. Let S be the set of elements of $\langle 3, 1, 5, 2 \rangle$ with this property,

$$S = \{(0, 1, 2, x_3, x_4) \in \langle 3, 1, 5, 2 \rangle \mid x_3 > 3\}$$

and note that the lattice paths associated with the elements of S are those which start by 2 horizontal steps, followed by a vertical step (cf. Figure 1). Now, $(0, 1, 2, x_3, x_4) \mapsto (0, x_3 - 3, x_4 - 3)$ defines a bijection between S and $\langle(3 + 1 + 5) - 3, 2\rangle$.

Consider the 5-composition $(3, 1, 5, 2, 4)$ of 15, define similarly to S the four sets $T \subseteq \langle 1, 5, 2, 4\rangle$, $U \subseteq \langle 5, 2, 4, 3\rangle$, $V \subseteq \langle 2, 4, 3, 1\rangle$ and $W \subseteq \langle 4, 3, 1, 5\rangle$. Then $|S| + |T| + |U| + |V| + |W| = \binom{6}{1} \binom{15}{7} + \binom{5}{1} \binom{10}{8} + \binom{8}{1} \binom{28}{10} + \binom{6}{1} \binom{15}{6} + \binom{5}{1} \binom{10}{9} = 27 + 30 + 52 + 21 + 35 = 165$. A similar construction for another 5-composition c of 15 gives again this number. If, e.g., $c = (2, 1, 7, 3, 2)$, we obtain in the same manner $\binom{7}{1} \binom{21}{9} + \binom{8}{1} \binom{28}{9} + \binom{9}{1} \binom{36}{10} + \binom{4}{1} \binom{6}{4} + \binom{2}{1} \binom{1}{8} = 42 + 44 + 54 + 10 + 15 = 165$.

The result that follows generalizes this situation.

Definition 4.5. Given integers t and $0 < r < k < n$, and a k -composition $\ell = (\ell_1, \dots, \ell_k) \in \mathbb{N}^k$ of n , let

$$s(\ell) := \sum_{i=0}^{k-1} \left| \langle (\sum_{j=1}^r \ell_{i+j}) + t, \ell_{i+r+1}, \dots, \ell_{i+k-1} \rangle \right|,$$

where indices are to be read modulo k .

Theorem 4.6. *The value of $s(\ell)$ does not depend on the k -composition ℓ .*

Proof. Note that, by definition, if we cyclically permute the elements of ℓ the value of $s(\ell)$ does not change. Hence, it is sufficient to prove that, given two k -compositions, $\mathbf{m} = (m_1, \dots, m_k)$ and $\ell = (\ell_1, \dots, \ell_k)$ such that

$$(m_1, m_2, \dots, m_{k-1}, m_k) = (\ell_1 - 1, \ell_2, \dots, \ell_{k-1}, \ell_k + 1),$$

we must have $s(\mathbf{m}) = s(\ell)$.

Let $s_i(\ell) = \left| \langle (\sum_{j=1}^r \ell_{i+j}) + t, \ell_{i+r+1}, \dots, \ell_{i+k-1} \rangle \right|$ and define $s_i(\mathbf{m})$ similarly. Then $s_0(\mathbf{m}) - s_0(\ell) = \langle L_r + t - 1, \ell_{r+1}, \dots, \ell_{k-1} \rangle - \langle L_r + t, \ell_{r+1}, \dots, \ell_{k-1} \rangle$, which is the negative of $\langle L_{r+1} + t - 1, \ell_{r+2}, \dots, \ell_{k-1} \rangle$ by Lemma 4.2.

In general, by subtracting and subsequently applying Lemma 4.2 term by term, we obtain

$$\begin{aligned} \sum_{i=0}^{k-1} (s_i(\mathbf{m}) - s_i(\ell)) &= \sum_{i=0}^{k-r} (s_i(\mathbf{m}) - s_i(\ell)) \\ &= - \left| \langle L_{r+1} + t - 1, \ell_{r+2}, \dots, \ell_{k-1} \rangle \right| \\ &\quad + \left| \langle (\sum_{j=1}^r \ell_{j+1}) + t, \ell_{r+2}, \dots, \ell_{k-1} \rangle \right| \\ &\quad + \left| \langle (\sum_{j=1}^r \ell_{j+2}) + t, \ell_{r+3}, \dots, \ell_{k-1} \rangle \right| \left| \langle \ell_1 - 1 \rangle \right| \\ &\quad \qquad \qquad \qquad \vdots \qquad \qquad \qquad \vdots \\ &\quad + \left| \langle (\sum_{j=1}^r \ell_{j+k-r-1}) + t \rangle \right| \left| \langle \ell_1 - 1, \ell_2, \dots, \ell_{k-r-2} \rangle \right| \\ &\quad + \left| \langle \ell_1 - 1, \ell_2, \dots, \ell_{k-r-1} \rangle \right| \end{aligned}$$

We prove that this number is zero by proving that the negative of the first summand, the size of $\mathfrak{X} = \langle L_{r+1} + t - 1, \ell_{r+2}, \dots, \ell_{k-1} \rangle$, is the sum of the other summands, each of which counts the elements with the same image by the function $f: \mathfrak{X} \rightarrow [k-r]$ such that

$$f(0, x_1, \dots, x_{k-r-1}) = \begin{cases} k - r, & \text{if } x_i < L_i, \forall i \leq k - r - 1; \\ \min\{i \mid x_i \geq L_i\}, & \text{otherwise.} \end{cases}$$

First, note that $f(X) = k - r$ if and only if $X \in \langle \ell_1 - 1, \ell_2, \dots, \ell_{k-r-1} \rangle$. If $X \notin \langle \ell_1 - 1, \ell_2, \dots, \ell_{k-r-1} \rangle$, then

$$\begin{aligned} f(X) \geq i &\iff \min\{i \mid x_i \geq L_i\} \geq i \\ &\iff \forall_{j < i}, x_j < L_j \end{aligned}$$

and hence

$$f(X) = i \iff (\forall_{j < i}, x_j < L_j) \wedge x_i \geq L_i.$$

Finally,

$$(0, x_1, \dots, x_{k-r-1}) \mapsto \left((0, x_1, \dots, x_{i-1}), (0, x_i - L_i + 1, \dots, x_{k-r-1} - L_i + 1) \right)$$

defines a bijection between $f^{-1}(\{i\}) \subseteq \mathfrak{X}$ and the set

$$\langle \ell_1 - 1, \ell_2, \dots, \ell_{i-1} \rangle \times \langle (\sum_{j=1}^r \ell_{j+i}) + t, \ell_{r+i+1}, \dots, \ell_{k-1} \rangle. \quad \square$$

4.2. Counting parking functions and rook words with a given type. Recall that a parking function of length n is a tuple $\mathbf{a} = (a_1, \dots, a_n) \in [n]^n$ such that the i th entry in ascending order is always at most $i \in [n]$. In other words,

$$\mathbf{a} \in \text{PF}_n \text{ if, for every } i \in [n], |\mathbf{a}^{-1}(\{i\})| \geq i.$$

Definition 4.7. Let $\mathbf{a} = (a_1, \dots, a_n) \in \mathbb{N}^n$ and suppose that $\{a_1, \dots, a_n\} = \{x_1, \dots, x_k\}$ with $x_i < x_j$ whenever $1 \leq i < j \leq k$.

The *reduced image* of \mathbf{a} is

$$\text{rim}(\mathbf{a}) = (x_1 - 1, \dots, x_k - 1) \in (\mathbb{N} \cup \{0\})^k;$$

the *coimage* of \mathbf{a} is the quotient set

$$\text{coim}(\mathbf{a}) = \{\bar{x} := \mathbf{a}^{-1}(\mathbf{a}(x)) \mid x \in [n]\}$$

ordered by

$$\bar{x} < \bar{y} \iff \mathbf{a}(x) < \mathbf{a}(y).$$

Let $\mathfrak{A} = (A_1, \dots, A_k)$ be an ordered (set) partition of $[n]$. The *length-vector* of \mathfrak{A} is

$$\ell(\mathfrak{A}) = (|A_1|, \dots, |A_{k-1}|).$$

In particular, the restriction of rim to $\text{coim}(\mathbf{a})$ is always injective and $\ell(\text{coim}(\mathbf{a})) = \text{rim}(\mathbf{a}) + (1, \dots, 1)$.

Lemma 4.8. Let \mathfrak{A} be an ordered partition of $[n]$ with length-vector $\ell(\mathfrak{A}) = (\ell_1, \dots, \ell_{k-1})$ and let $\mathbf{a}: [n] \rightarrow [n]$ be such that $\text{coim}(\mathbf{a}) = \mathfrak{A}$.

Then \mathbf{a} is a parking function if and only if

$$\text{rim}(\mathbf{a}) \in \langle \ell_1, \dots, \ell_{k-1} \rangle$$

and \mathbf{a} is a run r parking function if and only if

$$\begin{cases} \text{rim}(\mathbf{a}) = (0, 1, \dots, r-1, x_{r+1}, \dots, x_k) \\ (0, x_{r+1} - r, \dots, x_k - r) \in \langle (\sum_{i=1}^r \ell_i) - r, \ell_{r+1}, \dots, \ell_{k-1} \rangle \end{cases}$$

Proof. Follows immediately from the definitions. □

Note that $\mathbf{a} = (a_1, \dots, a_n) \in \text{RW}_n$ if and only if, for $i = a_1 - 1$, $\text{rim}(\mathbf{a})$ belongs to the set

$$\begin{aligned} \langle \underbrace{1, \dots, 1}_i, n - k + 1, \underbrace{1, \dots, 1}_{k-i-2} \rangle &= \\ &= \{(0, 1, \dots, i, x_{i+1}, \dots, x_{k-1}) \in [n-1]^k \mid i < x_{i+1} < \dots < x_{k-1} \leq n-1\}. \end{aligned}$$

Hence, if we denote by PF_n^r the set of run r parking functions of length n , for $\mathfrak{A} = (A_1, \dots, A_k)$, according to (4.3.1) and by definition

$$\begin{aligned} |\text{PF}_n \cap \text{coim}^{-1}(\mathfrak{A})| &= \det_{1 \leq i, j \leq k-1} \left(\binom{|A_1| + \dots + |A_i| - i + 1}{j - i + 1} \right) \\ |\text{PF}_n^r \cap \text{coim}^{-1}(\mathfrak{A})| &= \det_{r \leq i, j \leq k-1} \left(\binom{|A_1| + \dots + |A_i| - i}{j - i + 1} \right) \\ |\text{RW}_n \cap \text{coim}^{-1}(\mathfrak{A})| &= \binom{n-1-i}{k-1-i} \end{aligned}$$

Definition 4.9. Given an ordered partition \mathfrak{A} of $[n]$ and $\mathbf{a} \in [n]^n$, we say that \mathbf{a} is of type \mathfrak{A} if the coimage of \mathbf{a} is a cyclic permutation of \mathfrak{A} .

Theorem 4.10. Let $\mathfrak{A} = (A_1, \dots, A_k)$ be a partition of $[n]$ and let $1 \leq r \leq n$ for a natural number n . Then

- the number of **parking functions of type \mathfrak{A}** , as well as the number of **rook words of type \mathfrak{A}** , is $\binom{n}{k-1}$;
- the number of **run r parking functions of type \mathfrak{A}** , as well as the number of **run r rook words of type \mathfrak{A}** , is $r \binom{n-r-1}{k-r}$.

Proof. Consider the two ordered k -compositions of $[n]$, $\mathcal{C} = (|A_1|, |A_2|, \dots, |A_k|)$ and $\mathcal{D} = (n - k + 1, 1, \dots, 1)$, and apply Theorem 4.6 with different values of r and t .

For the first statement, take $r = 1$ and $t = 0$; in the notation thereof,

$$|\text{PF}_n \cap \text{coim}^{-1}(\mathfrak{A})| = s(\mathcal{C}) = \sum_{i=0}^{k-1} \binom{n-1-i}{k-1-i} = s(\mathcal{D}) = |\text{RW}_n \cap \text{coim}^{-1}(\mathfrak{A})|.$$

For the second statement, by taking $r = 1, \dots, k$ and $t = -r$ we obtain that

$$|\text{PF}_n^r \cap \text{coim}^{-1}(\mathfrak{A})| = s(\mathcal{D}) = r |\langle n - k, \underbrace{1, \dots, 1}_{k-r-1} \rangle| + 0,$$

since, for $\ell_1 = n - k + 1$ and $\ell_2 = \dots = \ell_k = 1$,

$$|\langle (\sum_{j=1}^r \ell_{i+j}) - r, \ell_{i+r+1}, \dots, \ell_{i+k-1} \rangle| = \begin{cases} |\langle n - k, 1, \dots, 1 \rangle|, & \text{if } i \in [r]; \\ 0, & \text{otherwise.} \end{cases}$$

This shows that the number of run r parking functions of type \mathfrak{A} is $r \binom{n-r-1}{k-r}$, since

$$\langle n - k, 1, \dots, 1 \rangle = \{(x_1, \dots, x_{k-r}) \mid 0 < x_1 < \dots < x_{k-r} \leq n - r - 1\}.$$

Finally, note that, for example, all the type \mathfrak{A} elements $\mathbf{a} = (a_1, a_2, \dots, a_n) \in [n]^n$ with $a_1 = 1$ share the same coimage, and that there are $\binom{n-r-1}{k-r}$ such rook words with run r ,

since they are determined by the last $k - r$ strictly increasing coordinates of $\text{rim}(\mathbf{a})$, all of them greater than r and less than n . The same happens if $a_1 = i$ for $1 \leq i \leq r$, and a_i cannot be greater than r , by definition. \square

We note that the first part of Theorem 4.10 can be obtained directly from [10, Cyclic Lemma], where the following bijection is defined. Let $\mathbf{b} = \mathbf{a}$ if $\mathbf{a} \in \text{PF}_n \cap \text{RW}_n$ and, if $\mathbf{a} \in \text{PF}_n \setminus \text{RW}_n$ and $m = \max([a_1] \setminus \mathbf{a}([n]))$, let $\mathbf{b} = (b_1, \dots, b_n) \in [n]^n$ be such that

$$a_i \equiv b_i + m \pmod{n};$$

then $\mathbf{a} \mapsto \mathbf{b}$ defines a bijection between the set of parking functions and the set of rook words of a given type.

Theorem 4.11. *For every $n \in \mathbb{N}$,*

$$\mathcal{RPF}_n(t) = \mathcal{RRW}_n(t).$$

Proof. Follows immediately from Theorem 4.10. \square

5. COUNTING ROOK WORDS WITH A GIVEN RUN

Given positive integers n and r such that $r \leq n$, let

$$\text{RW}_n^r = \{f \in \text{RW}_n \mid \text{run}(f) = r\}$$

and for $\mathbf{a} = (a_1, \dots, a_n) \in \text{RW}_n^r$ let

$$\mathbf{r} = \mathbf{r}(\mathbf{a}) = (i_1, \dots, i_r)$$

where $i_j = \min\{i \in [n] \mid a_i = j\}$ for every $j \in [r]$ (compare with the definition of Run in page 6).

Theorem 5.1. *For every integers $1 \leq r \leq n$,*

$$(5.1.1) \quad [t^r](\mathcal{LLT}_n(t)) = r! \sum_{i_1 + \dots + i_r = n-r} (n-1)^{i_1} (n-2)^{i_2} \dots (n-r)^{i_r}$$

$$(5.1.2) \quad = r \sum_{j=0}^{r-1} (-1)^j \binom{r-1}{j} (n-1-j)^{n-1}.$$

Proof. We have seen before that

$$[t^r](\mathcal{LLT}_n(t)) = [t^r](\mathcal{RRW}_n(t)) = |\text{RW}_n^r|.$$

Given $\mathbf{a} \in \text{RW}_n^r$ and $\pi \in \mathfrak{S}_r$, let $\pi\mathbf{a}$ be the element of $[n]^n$ defined by

$$(\pi\mathbf{a})(j) = \begin{cases} \pi(a_j) & \text{if } a_j \leq r; \\ a_j & \text{if } j > r. \end{cases}$$

Note that $\pi\mathbf{a} \in \text{RW}_n^r$ if and only if $\mathbf{a} \in \text{RW}_n^r$. Owing to this, the left-hand side of (5.1.1) is equal to $r!$ times the number of elements of

$$A = \{\mathbf{a} \in \text{RW}_n^r \mid \mathbf{r}(\mathbf{a}) = (i_1, \dots, i_r) \text{ with } 1 = i_1 < i_2 < \dots < i_r\}.$$

Now, $\mathbf{a} = (a_1, \dots, a_n) \in A$ if and only if, for every $1 \leq \ell \leq n$,

- $a_\ell \notin \{i_{j+1}, \dots, i_r, r+1\}$, if $i_j < \ell < i_{j+1}$ for some $j \in [r-1]$;
- $a_\ell \neq r+1$, if $\ell > i_r$.

This gives (5.1.1) for $i_{r+1-j} = i_{j+1} - i_j - 1$ and $i_1 = n - i_r$, respectively.

We note that the right-hand side of (5.1.2) is, by the Inclusion-Exclusion Principle, r times the number of elements of

$$B = \{f: [n-1] \rightarrow [n-1] \mid [r-1] \subseteq f([n-1])\}.$$

Given $\ell \in [n]$ with $\ell \leq r < n$, consider the bijection $\varphi_\ell: [n] \setminus \{r+1\} \rightarrow [n-1]$ such that

$$\varphi_\ell(j) = \begin{cases} j, & \text{if } j < \ell; \\ r, & \text{if } j = \ell; \\ j-1, & \text{if } j > \ell. \end{cases}$$

and note that $[r-1] \subseteq \varphi_\ell([r])$. Now, $F(a_1, \dots, a_n) = (a_1, \varphi_{a_1}(a_2), \dots, \varphi_{a_1}(a_r))$ clearly defines a bijection from RW_n^r to $[r] \times B$. \square

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